# Using BOINC to enumerate mutually orthogonal Latin squares 

Gerdus Benadé, Alewyn Burger, Jan van Vuuren

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## Latin squares

## Definition (Latin square)

A Latin square of order $n$ is an $n \times n$ square where every cell contains one of the symbols in the set $\{0,1, \ldots, n-1\}$ such that no symbol is repeated in any row or column.

| 2 | 1 | 0 | 3 | 2 | 0 | 3 | 1 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 3 | 2 | 1 | 1 | 3 | 0 | 2 |
| 3 | 0 | 1 | 2 | 0 | 2 | 1 | 3 |
| 1 | 2 | 3 | 0 | 3 | 1 | 2 | 0 |

## Latin squares - Orthogonality

## Definition

Two latin squares of the same order are orthogonal if, when superimposed, each of the possible $n^{2}$ ordered pairs occur exactly once.

| 2 | 1 | 0 | 3 |  | 2 | 0 | 3 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 3 | 2 | 1 | 1 |  |  |  |
| 3 | 0 | 1 | 2 | 1 | 3 | 0 | 2 |
| 1 | 2 | 3 | 0 |  | 3 | 2 | 1 |

## Latin squares - Orthogonality

## Definition

A set $\left\{L_{1}, \ldots, L_{k}\right\}$ of $k \geq 2$ latin squares of order $n$ is orthogonal if any two distinct latin squares are orthogonal. We call this a set of $k$ mututally orthogonal latin squares ( $k$-MOLS) of order $n$.

| 2 | 1 | 0 | 3 |
| :--- | :--- | :--- | :--- |
| 0 | 3 | 2 | 1 |
| 3 | 0 | 1 | 2 |
| 1 | 2 | 3 | 0 |


| 2 | 0 | 3 | 1 |
| :--- | :--- | :--- | :--- |
| 1 | 3 | 0 | 2 |
| 0 | 2 | 1 | 3 |
| 3 | 1 | 2 | 0 |


| 0 | 1 | 2 | 3 |
| :--- | :--- | :--- | :--- |
| 1 | 0 | 3 | 2 |
| 2 | 3 | 0 | 1 |
| 3 | 2 | 1 | 0 |

## Permutations and main classes

Permuting/re-arranging the rows, symbols and columns of a MOLS give a structurally similar MOLS.

## Definition (Main class)

A MOLS, together with all its permutations, is called a main class.

## Definition (Class representative)

Every main class has a lexicographical smallest element, called the class representative.

## Enumerating main classes of MOLS

## Problem

How many main classes of $k-M O L S$ of order $n$ are there?
Approach: Count the class representatives


## Enumerating main classes of MOLS

## Problem

How many main classes of $k-M O L S$ of order $n$ are there?

| $n$ | $k$ |  |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :---: |
|  | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |  |
| 3 | 1 |  |  |  |  |  |  |  |  |
| 4 | 1 | 1 |  |  |  |  |  |  |  |
| 5 | 1 | 1 | 1 |  |  |  |  |  |  |
| 6 | 0 | 0 | 0 | 0 |  |  |  |  |  |
| 7 | 7 | 1 | 1 | 1 | 1 |  |  |  |  |
| 8 | 2165 | 39 | 1 | 1 | 1 | 1 |  |  |  |
| 9 | $\geq 1$ | $\geq 1$ | $\geq 1$ | $\geq 1$ | $\geq 1$ | $\geq 1$ | 19 |  |  |
| 10 | $\geq 1$ | $?$ | $?$ | $?$ | $?$ | 0 | 0 | 0 |  |

Table: The number of structurally different $k$-MOLS of order $n$ for $n \in\{3,4, \ldots, 10\}$.

## Enumeration algorithm - Universals


$<0,1,2,3>$

$<0,2,3,1>$

## Enumeration algorithm - Universals


$<0,1,2,3>$
$<1,0,3,2>$

$<0,2,3,1>$

A partial MOLS needs certain properties.
$\mathbb{P}=\left\{\begin{array}{l}\text { Every square in } M \text { has to be 'Latin' } \\ \text { All pairs in } M \text { orthogonal } \\ \text { No 'smaller' partial MOLS in this main class }\end{array}\right.$

## Enumeration algorithm - Example



## Enumeration algorithm - Example



## Enumeration algorithm - Example



## Enumeration algorithm - Example



## Enumeration algorithm - Example



## Enumeration algorithm - Branches

Table: The number of branches on after every symbol for a 3-MOLS of order $n$.

|  | After universal |  |  |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $n$ | 1 | 2 | 3 | 4 | 5 | 6 |
| 3 | 1 | 1 | 1 |  |  |  |
| 4 | 1 | 1 | 1 | 1 |  |  |
| 5 | 2 | 4 | 2 | 2 | 1 |  |
| 6 | 3 | 20 | 0 | 0 | 0 | 0 |
| 7 | 14 | 10529 | 3800 | 3 | 3 | 3 |
| 8 | 45 | 15948763 | 1546241258 | 18877734 | 216 | 168 |
| 9 | 269 | $2.89 \times 10^{10}$ | $8.48 \times 10^{14}$ | $2.68 \times 10^{15}$ |  |  |
| 10 | 1700 | $1.21 \times 10^{14}$ | $2.42 \times 10^{21}$ | - |  |  |

The number of nodes in the tree grows very quickly, making the complete enumeration for orders 9 and up difficult, if not impossible.

## What next?

| $n$ | $k$ |  |  |  |  |  |  |  |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
|  | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |
| 3 | 0 |  |  |  |  |  |  |  |
| 4 | 0 | 0 |  |  |  |  |  |  |
| 5 | 0 | 0 | 0 |  |  |  |  |  |
| 6 | 0 | 0 | 0 | 0 |  |  |  |  |
| 7 | 4 | 6 | 2 | 1 | 1 |  |  |  |
| 8 | $8 d$ | $9 d$ | $3 h$ | $1 h$ | $650 s$ | $113 s$ |  |  |
| 9 |  | $\approx 14 y$ |  |  |  |  |  |  |
| 10 |  | $\approx 60 y$ |  |  |  |  |  |  |

- Fix graphics, portability, GPU
- Test within our university network
- Public launch if feasible


## Natural partition



## Job size estimation

- Job size range from 1 second to 90000 (3-MOLS order 8) hard to provide accurate estimate and maximum fpops
- Workunits are too long, but there are too many to generate them on a deeper level. Enumerate partial subtrees using checkpoints as new workunits.

Starting points


## Interactivity



## Interactivity



Keep very little work on hand, use visualization to generate workunits on demand (suitable?).
Alternative: "Poll" to decide which area to explore next.
Should work best with shorter workunits, see the consequences of your choices

